`

**DotNetVault**

Synchronization Library and Static Analyzer for C# 8.0

Project Description

Author: Christopher P. Susie

Copyright © 2019, CJM Screws, LLC. All rights reserved.

**Table of Contents**

[1. Introduction. 5](#_Toc28769253)

[a. Currently used synchronization methods 5](#_Toc28769254)

[b. Problems with current lock-based mechanisms 6](#_Toc28769255)

[i. Primary problem with current mechanisms is they protect data only when programmers follows convention; Also, try … finally syntax error prone 6](#_Toc28769256)

[ii. Atomic operations are highly useful alternative but not easy to understand and scope of usefulness limited compared to locks 7](#_Toc28769257)

[iii. C#’s lock mechanism is not timed when used in its RAII form and is bug prone when used in its try…finally form 7](#_Toc28769258)

[iv. C#’s Monitor Lock mechanism is recursive 7](#_Toc28769259)

[v. Carefully crafted objects can be the most effective solution, but these are often not possible or maintainable by all given changing requirements 11](#_Toc28769260)

[c. DotNetVault isolates protected data and prevents access to it without first obtaining a lock 11](#_Toc28769261)

[i. C# 8’s Disposable ref struct is used to isolate obtained locks on the stack and ensure prompt release in all cases 12](#_Toc28769262)

[ii. Static analysis prevents leakage or mingling of shared mutable state 13](#_Toc28769263)

[2. Prerequisites 14](#_Toc28769264)

[3. Installation 14](#_Toc28769265)

[4. Usage Guide 14](#_Toc28769266)

[a. Concept of Vault-Safety 14](#_Toc28769267)

[b. Overview of Tools 16](#_Toc28769268)

[i. *Vaults* 16](#_Toc28769269)

[ii. *LockedResources* 16](#_Toc28769270)

[c. Vaults In-Depth 18](#_Toc28769271)

[i. Functionality Common to All Vaults (intended for public consumption) 18](#_Toc28769272)

[ii. *BasicVault<T>* 21](#_Toc28769273)

[iii. *MutableResourceVault<T>* 22](#_Toc28769274)

[iv. *CustomizableMutableResourceVault<T>* 25](#_Toc28769275)

[d. LockedResourceObjects In-Depth 25](#_Toc28769276)

[i. Common Functionality 25](#_Toc28769277)

[*ii.* *LockedVaultObject<TVault, [VaultSafeTypeParam] T>* 26](#_Toc28769278)

[*iii.* *LockedVaultMutableResource<TVault, TResource>* 28](#_Toc28769279)

[5. Static Analyzer Rules 37](#_Toc28769280)

[a. DotNetVault\_UsingMandatory 37](#_Toc28769281)

[i. Summary 37](#_Toc28769282)

[ii. Known Bug (#50) 37](#_Toc28769283)

[b. DotNetVault\_VaultSafe 38](#_Toc28769284)

[c. DotNetVault\_VsDelegateCapture 39](#_Toc28769285)

[6. Attributes 42](#_Toc28769286)

[a. VaultSafeAttribute 42](#_Toc28769287)

[b. UsingMandatoryAttribute 43](#_Toc28769288)

[c. VaultSafeTypeParamAttribute 43](#_Toc28769289)

[d. NoNonVsCaptureAttribute 43](#_Toc28769290)

[e. NotVsProtectableAttribute 44](#_Toc28769291)

[7. Licensing 44](#_Toc28769292)

[a. Software License 44](#_Toc28769293)

[b. Documentation License 45](#_Toc28769294)

**Table of Figures**

[Figure 1 – RAII style Lock and Try…Finally Equivalent 6](#_Toc28769295)

[Figure 2 -- Typical Wrapper Around Well-Known Interface 8](#_Toc28769296)

[Figure 3 – Typical Flawed Attempt at Post Hoc Thread Safety 9](#_Toc28769297)

[Figure 4 – Demonstrates Lack of State Consistency Between Calls. 10](#_Toc28769298)

[Figure 5 – Typical Mechanism for Achieving Thread Safety After the Fact 10](#_Toc28769299)

[Figure 6 – Public Properties Common to All Vaults 18](#_Toc28769300)

[Figure 7 – Public Methods Common to All Vaults 19](#_Toc28769301)

[Figure 8 – Lock and Spinlock Overloads 21](#_Toc28769302)

[Figure 9 – Correct and Incorrect Creation of MutableResourceVault 23](#_Toc28769303)

[Figure 10 – More Elaborate Correct and Incorrect Creation of MutableResourceVault 24](#_Toc28769304)

[Figure 11 – Care is Needed When Working With Mutable Structs 27](#_Toc28769305)

[Figure 12 – Output Mutable Struct Demonstration 28](#_Toc28769306)

[Figure 13 -- Special Delegates Used By LockedVaultMutableResource Objects to Prevent Leakage and Mingling of State 29](#_Toc28769307)

[Figure 14 -- (cont’d) Special Delegates Used By LockedVaultMutableResource Objects to Prevent Leakage and Mingling of State 30](#_Toc28769308)

[Figure 15 -- VaultQuery Demonstration 32](#_Toc28769309)

[Figure 16 -- VaultQuery Demo Output 32](file:///G:\CjmScrews\Source\VaultAnalyzer\MasterBranch_Epic36_Story40_Tasks41_42_43_44_45\DotNetVault%20Description.docx#_Toc28769310)

[Figure 17 -- VaultAction Demonstration 33](#_Toc28769311)

[Figure 18 -- VaultAction Demo Output 33](file:///G:\CjmScrews\Source\VaultAnalyzer\MasterBranch_Epic36_Story40_Tasks41_42_43_44_45\DotNetVault%20Description.docx#_Toc28769312)

[Figure 19 – VaultMixedOperation Demonstration 34](#_Toc28769313)

[Figure 20 -- VaultMixedOperation Demo Output 35](file:///G:\CjmScrews\Source\VaultAnalyzer\MasterBranch_Epic36_Story40_Tasks41_42_43_44_45\DotNetVault%20Description.docx#_Toc28769314)

[Figure 21 – Demonstration of Extension Methods to Simplify Usage 36](#_Toc28769315)

[Figure 22 -- Output of Extension Method Demo 36](file:///G:\CjmScrews\Source\VaultAnalyzer\MasterBranch_Epic36_Story40_Tasks41_42_43_44_45\DotNetVault%20Description.docx#_Toc28769316)

[Figure 23 -- Bug #50 Demonstration 38](#_Toc28769317)

[Figure 24 -- Vault-Safe Convenience Wrappers 40](#_Toc28769318)

[Figure 25 – Usage of Vs Convenience Wrappers 41](#_Toc28769319)

[Figure 26 -- Usage Wrapper Demo Output 42](#_Toc28769320)

[Figure 27 -- Contents of Whitelist.txt 42](#_Toc28769321)

[Figure 28-- Contents of condit\_generic\_whitelist.txt 43](file:///G:\CjmScrews\Source\VaultAnalyzer\MasterBranch_Epic36_Story40_Tasks41_42_43_44_45\DotNetVault%20Description.docx#_Toc28769322)

# Introduction.

DotNetVault is a library and static code analysis tool that makes managing shared mutable state in multi-threaded applications more manageable and less error prone. Where errors do still occur, they are easier to locate and identify.

Managing shared mutable state in multithreaded environments is one of the most difficult feats to accomplish in complex application development. The tools available generally rely on each programmer’s knowledge of the synchronization mechanisms employed. Each programmer in the project must know and follow them always, even when that programmer’s task or expertise lies outside threading issues. Even highly experienced multi-threading programmers make mistakes because focus on the domain area or a difficult unrelated task may divert his or her attention. These mistakes are difficult to reproduce, diagnose, find and correct. They rank among the costliest bugs in software projects.

## Currently used synchronization methods

One way to eliminate issues of synchronization is to eliminate the use of shared mutable state. If all state is mutable but not shared or shared but not mutable, worries about thread synchronization disappear. For most application development, mutability of shared state can be *reduced* but rarely *eliminated*. Furthermore, the ubiquity of reference types in C# works to thwart attempts at preventing sharing. Thus, in C#, programmers are frequently left with traditional synchronization primitives that require a portion of the engineer’s mental focus, leaving her unable to focus entirely on solving the problem at hand. The most frequently used primitive in C# to synchronize state is the *lock* statement.[[1]](#footnote-1) Lock is a simple RAII[[2]](#footnote-2)-like mechanism around a more convoluted and bug-prone syntax involving *try … finally*:



Figure [[3]](#footnote-3) – RAII style Lock and Try…Finally Equivalent

There are also other similar synchronization mechanisms available to the C# programmer: The *System.Mutex*[[4]](#footnote-4) class, *ReaderWriterLockSlim*[[5]](#footnote-5)*, SpinLock*[[6]](#footnote-6)and many others. In addition to these locking mechanisms, there are lock-free thread safety mechanisms that provide thread-safety through atomic operations.[[7]](#footnote-7)

## Problems with current lock-based mechanisms

### Primary problem with current mechanisms is they protect data only when programmers follows convention; Also, try … finally syntax error prone

Of the lock-based synchronization mechanisms provided, there are many drawbacks. The primary drawback is that they rely on programmer discipline to only access the protected objects when the lock is obtained. A related drawback is that, except for *lock* which incorporates guaranteed scope-based release of the lock, the programmer must remember to release the lock even in exceptional cases, but not release it too soon. Most, if not all, of the synchronization mechanisms other than *lock* do not provide an RAII-like way to do this, so try finally must be employed, which is a convoluted syntax, susceptible to being forgotten or mistakenly implemented.[[8]](#footnote-8)

### Atomic operations are highly useful alternative but not easy to understand and scope of usefulness limited compared to locks

Furthermore, atomic operations, unless you are dealing with the limited number of types for which atomic operations exist, are no panacea: they are frequently difficult to understand for programmers less experienced with multithreading and offer a much smaller window of atomicity than that provided by the mutex/lock-based synchronization mechanisms. Nevertheless, they do provide a valuable alternative to mutexes when used judiciously. DotNetVault makes extensive use of them.

### C#’s lock mechanism is not timed when used in its RAII form and is bug prone when used in its try…finally form

The heavy onus on the programmer to refrain from using protected resources (and knowing which resources are protected and which protected resources are associated with which mutex) is far from the only other drawback of the commonly employed mutexes. When multiple mutexes are held simultaneously, the user must be especially disciplined to make sure that *every time* they are acquired, the same relative ordering is applied. Failure to do this can sometimes result in silent deadlock of the program but may not be reliably reproducible. One way to get around this possibility is to set a timeout when you attempt to acquire this mutex: instead of silent deadlock an exception will be thrown, making it easy to identify the problem and correct it. Unfortunately, C# does not provide an out-of-the-box RAII-based mechanism to add a time limit, forcing you to rely on convoluted try-finally constructs as shown above.[[9]](#footnote-9)

### C#’s Monitor Lock mechanism is recursive

Another problematic consideration is C#’s usage of recursive mutexes (and similar locking mechanisms). Most mutex mechanisms in C# are recursive, recursive by default or at least optionally recursive. The most used mechanism is the *Monitor.TryEnter* method or its more reliable RAII *lock* shorthand.[[10]](#footnote-10) This type of mutex is, by default, recursive. The primary desirable use-case for a recursive mutex is the adaptation of code that was not designed to be thread-safe to some semblance of thread-safety, *without changing the API*.[[11]](#footnote-11) This usage, however, should be considered a last resort.[[12]](#footnote-12)

#### Recursive mutexes are intended to be used to adapt non-thread safe objects to a thread-safe context

To see how the recursive mutex can work to adapt non-thread-safe code to thread-safety while remaining problematic, consider the following custom list wrapper that presumably would add some domain-specific functionality to a standard List<T>:



Figure -- Typical Wrapper Around Well-Known Interface

As frequently happens, later in the development of the software, it is discovered or decided that a thread-safe version of this type is required. The following code results is a typical result:



Figure – Typical Flawed Attempt at Post Hoc Thread Safety

#### While it may achieve this on a syntactic level, it usually fails to do so on a semantic level leading to unmaintainable code

Note that the *GetEnumerator*() method returns an enumerator to a snapshot copy of the collection rather than an enumerator to the actual collection.[[13]](#footnote-13) Also, the synchronization object is exposed (thus adding to the API but not removing from it or altering its signatures). The reason for the exposure of the synchronization object is shown in this example:



Figure – Demonstrates Lack of State Consistency Between Calls.

The problem here should be immediately apparent: the lock is released after execution of *IndexOf* is complete then is ***reacquired*** for the call to the set accessor of the indexer property in *list[idxToFirstJaneInList].* The list’s contents may change between the two calls. So, given the recursive nature of *lock*s in C# and the handy public *SyncObject* property exposed by the thread-safe wrapper, the solution presents itself:



Figure – Typical Mechanism for Achieving Thread Safety After the Fact

Notice that in attempting to preserve the same API for the object, we have really introduced a vexing problem that will propagate to all client code. It is rare that you only do one “operation” on a list: typically, you do multiple operations on a list and each successive operation relies on the previous operations’ results still being accurate. Thus, nearly everywhere your new thread-safe implementation is used, you will have to make sure that you obtain a recursive lock. Also, it will not be a compiler error or obviously wrong on inspection of any code block, *as a unit*, if this lock is not first obtained. Recursively obtaining mutexes is not free from a performance penalty compared to non-recursive mutexes.

This approach fails: while it *appears* to preserve the same API the thread-unsafe List had, it does so *only with syntax*. To make the thread-safe version equivalent to the original version *semantically*, additional actions must be taken by the caller, violating the Liskov substitution principle[[14]](#footnote-14) and leading to bugs and unmaintainable code.

### Carefully crafted objects can be the most effective solution, but these are often not possible or maintainable by all given changing requirements

One approach superior to the recursive mutex approach is to design a special purpose object to access the list. While this may be superior *in theory*, it requires you to anticipate which sequence of access to the list will be required and to expose them all in the object’s public interface. This is more difficult than it sounds in the light of changing requirements and the differing skill levels of project contributors who may be forced to alter the thread safe object as new needs are discovered.

## DotNetVault isolates protected data and prevents access to it without first obtaining a lock

DotNetVault takes its inspiration from the synchronization mechanisms provided by Rust language and the Facebook Folly C++ synchronization library. These synchronization mechanisms observe that the mutex should *own* the data they protect. You literally cannot access the protected data without first obtaining the lock. RAII destroys the lock when it goes out of scope – even if an exception is thrown or early return taken.[[15]](#footnote-15)

In a sense, languages like C++ and Rust that heavily use types with value semantics, incorporate the concept of *ownership* of a resource and provide *move* semantics are more easily able to provide isolation of resources than a language like C# where so many we access so many objects through garbage collected references: these references, thanks to garbage collection, can be shared promiscuously and by their very nature almost beg to be shared. While garbage collection saves us from the nightmare of keeping track of who needs to free all these freely shared references to heap objects[[16]](#footnote-16), it works at cross-purposes to isolation aimed at thread-safety.

Two factors have made an isolating approach to protected resources more obtainable in C#: one very recent and one somewhat recent: disposable ref structs and Roslyn analyzers. DotNetVault uses the C# 8.0 language feature (disposable ref structs) and specifically tailored static analysis tools to isolate protected data and force locks to them to be acquired and freed in accordance with RAII.

### C# 8’s Disposable ref struct is used to isolate obtained locks on the stack and ensure prompt release in all cases

A *ref struct* is a value type that can *only* reside on the stack. Structs in general (along with *enums*) in C# are value types. Some people mistakenly believe that value types “live on the stack.” This notion is deeply flawed.[[17]](#footnote-17) A struct variable contains its value within itself. It can end up on the heap under two primary circumstances:

1. It is a field in a reference type
2. It is an item in a generic array (*T*[])
3. It is boxed[[18]](#footnote-18)

For example, if you have a local variable of struct type and assign it to another local variable of type object (or, for example, *IDisposable*, if the struct implements that interface), then the struct is copied to the heap and a pointer to that boxed struct (and v-table) is stored in your object/*IDisposable* variable. *Ref structs,* with their counterintuitive name, are *truly impossible to store on the heap or in static memory*. They cannot be a field in a class or ordinary struct, they cannot be assigned to an object, they cannot be stored in static variables, they cannot be boxed.[[19]](#footnote-19)

Ref structs have been with us since C# 7.2. Note that since they cannot be boxed, they also cannot implement interfaces, including IDisposable. *IDisposable* is a feature that allows, together with a using statement, RAII-like semantics to be obtained in C#. A variable declared in a using statement is guaranteed to call its *Dispose* method when its scope ends (regardless of exception or early return). With C# 8.0, you may use “using” statement (or now, declaration) with a ref struct has a void *Dispose*() method even though it cannot implement the *IDisposable* interface.

This gives us an object that cannot be stored on the heap in any way: it cannot be boxed, it cannot be captured in a closure, it cannot be stored in a class or ordinary struct field. It cannot be an element in an ordinary array. Its lifetime will be no longer than its scope, by inexorable guarantee. DotNetVault exploits the restricted deterministic lifetime of these types of objects to assist it in isolating a protected resource and ensuring it is properly freed in short order and not stored somewhere for an indefinite period.

### Static analysis prevents leakage or mingling of shared mutable state

By itself, this would not be enough to provide guarantees of isolation. Roslyn analyzers provide a framework where domain specific rules can be imposed on a program and enforced both by IntelliSense and at Build Time. DotNetVault has developed rules that enforce the isolation of resources within a vault and ensure that shared mutable state cannot leak out of a protected resource nor can external shared mutable state be accidentally mingled with it. It also can impose a rule that requires the use of a using statement or declaration by the immediate caller of a function returning a ref struct, making it a compiler error to neglect this.

In summary, DotNetVault provides the following advantages to managing shared mutable state in multithreaded applications:

1. The protected resource cannot be accessed before a lock has been obtained or after it has been released
2. The lifetime of the lock obtained cannot exceed the function in which it is called
3. The lock will be disposed of properly even in face of an exception or early return with nothing more onerous that a using declaration
4. When the lock is obtained, the user will find it very difficult to leak shared mutable state out of the lock object – most attempts will be thwarted with compiler errors: the goal is 100% isolation, but, realistically, making it very difficult to use incorrectly is probably the best that can be hoped for
5. The mechanism used is not recursive, but it will not deadlock unless a “wait forever” decision is explicitly made by the caller: by default, all attempts to obtain the lock, including in a recursive scenario, will timeout.

# Prerequisites

This library and static analyzer target the C# language. It requires features that are available only in C# 8.0 and later. Generally, this will require .NET Core 3.0 or later. .NET Framework 4.8 does not, out-of-the-box, support all C# 8 features. It is possible, however, to use this library and analyzer with .NET 4.8 or .NET Standard 2.0 provided you set the language version to 8.0 manually in the *csproj* file by adding the following line in a property group that applies to all build configurations-

<LangVersion>8.0</LangVersion>

The features of C# 8.0 used by this project are “syntax-only” features[[20]](#footnote-20) that work equally well with .NET Standard 2.0, .NET 4.8 and .NET Core 3.0 and perhaps other implementations. Indeed, this project was written and tested using .NET Standard 2.0 with C# 8.0 manually enabled as shown above. If using .NET Standard 2.0 or .NET Framework 4.8, you should avoid attempting to use features of C# 8.0 that require runtime/framework support.

# Installation

Installation is performed by using NuGet to install the package. Officially, Visual Studio 2019 is supported. Support for installation into current and future versions of JetBrains Rider is under consideration.

# Usage Guide

## Concept of Vault-Safety

Vault-Safety is a concept used extensively by this library and static analyzer. Vault-Safety is a characteristic of certain types that makes it easier to protect them leaking outside of the protection of the *Vault* and *Lock* objects used in this software. If a type is Vault-Safe, it is easier to ensure that protected data will not have references that can mutate its state or sneak out from the *Vault* barrier.

A type is Vault-Safe if and only if it has one or more of the following characteristics:

1. It is an *unmanaged[[21]](#footnote-21)* value type – that is a value type that contains no reference types (or unsafe pointers or dynamic types) at any level of recursion in its object map. Any type that meets the *unmanaged* type constraint is *automatically* considered Vault-Safe.
2. It is a concrete reference type that is *sealed* and contains only immutable data. The string type is a paradigmatic example of such a type. It is always considered Vault-Safe by the static analyzer. Also considered Vault-Safe are the collections in the *System.Collections.Immutable* namespace (except their builders), but **only to the extent that all their type arguments are vault-safe**.
3. It is a value type that only contains fields (static and member) which themselves contain only:
   1. Other value types and/or
   2. Reference types that meet the criteria of #ii, *supra* (no mutable state and sealed).

The reason that such types are easy to isolate is simple: assignment of those types results in either an *actual deep copy* of the object (for *unmanaged* types) or, for other vault-safe types, *effectively a deep copy*. Deep copies of data, whether true deep copies or quasi deep copies due to immutability, cannot be used to change the value protected. Concerns about references to the protected data sneaking out of the ambit of their protective isolation do not apply to these types.

*Unmanaged* types require no special annotation to be considered vault-safe: they are manifestly vault-safe by their very nature. For reference types and value types that do not comply with the *unmanaged* constraint, you must either annotate the type with the *VaultSafeAttribute[[22]](#footnote-22)* or, alternatively place its type in the whitelist file[[23]](#footnote-23). It is recommended that the *VaultSafeAttribute* be employed for types over which you have control and the whitelist for types beyond your ability to annotate with attributes (for example, because they are part of a 3rd-party API).

## Overview of Tools

As a consumer of this project, there are two ready-made categories of types provided for your use. *Vaults* isolate protected resources and prevent concurrent access to them or any values in their object graph. To obtain temporary access to an object protected by a *Vault* you call the one of the *Vault’s Lock()* or *SpinLock()* methods to obtain the second category of type: the *LockedResource*. The *LockedResource* is an RAII type that may only be stored on the stack and must be disposed by the same Method (or Property) that obtained it. In fact, the method or property that obtains it must guard it with a *using* statement or declaration. Failing to do so is a compiler error.

### *Vaults*

Vault objects all inherit from the abstract base class *Vault<T>*. There are two Vaults provided for your use out-of-the-box: the *MutableResourceVault<T>* and the *BasicVault<T>*. The difference between which should be used depends on whether *T* is Vault-Safe*.[[24]](#footnote-24)* The *BasicVault<T>* should be used to protect Vault-Saferesources because the *LockedResource* provided by the *BasicVault<T>* is minimally restrictive*.* The *MutableResourceVault<T>* should be used to protect resources that are not Vault-Safe: the *LockedResource* object it provides when you obtain a locked resource prevents references to the protected data from leaking out and prevents non-vault safe external data from creeping in.[[25]](#footnote-25) Finally an abstract *CustomizableMutableResourceVault<T>*, inherits from *MutableResourceVault<T>* and exits to allow you to make your own version of the *MutableResourceVault<T>* with your own *LockedResource* type to provide an accessible API for frequently used operations on a type that is not Vault-Safe.

### *LockedResources*

Every publicly accessible *Vault* object contains public methods called *Lock* and *SpinLock*. These methods, all of which are annotated with the *UsingMandatory* attribute, “check-out” the protected resource give you temporary access to the resource mediated by a LockedResourceobject. LockedResource objects are Disposable *ref structs*[[26]](#footnote-26)that grant you limited access to the resource and automatically return the resource to the owning *Vault* when their lifetime ends. BasicVault<T> provides a LockedResource of type *LockedVaultObject<TVault, T>* and *MutableResourceVault<T>* provides a LockedResource of type *LockedVaultMutableResource<TVault, TResource>*. *LockedVaultObject<TVault, T>* simply exposes the protected resource to you as a read/write Property: this is all the protection a Vault-Safe resource requires. The *LockedMutableResourceVault<TVault, TResource>* restricts your interaction with the protected resource to the use of certain specially annotated delegates. The annotations on these delegates cause the static analyzer to refuse to compile your code if it cannot prove the code you passed in the delegate will not leak mutable state out of the resource or permit externally accessible mutable state from becoming part of the resource.

## Vaults In-Depth

### Functionality Common to All Vaults (intended for public consumption)

#### Public Read Only Properties:

|  |  |  |
| --- | --- | --- |
| Type | Name | Description |
| bool | DisposeInProgress | The dispose operation is currently in progress |
| bool | IsDisposed | The vault has been disposed. |
| TimeSpan | DefaultTimeout | The default maximum amount of time to spend attempting to obtain a lock. |
| TimeSpan | DisposeTimeout | How long to wait for a locked resource to be returned to the vault when the Dispose method is called. If the locked resource is not returned to the vault in time, an exception will be thrown. |
| TimeSpan | SleepInterval | How long should the thread sleep between failed attempts to obtain a lock. Typically, a few milliseconds. Not used at all when *SpinLock* is called. |

Figure – Public Properties Common to All Vaults

#### Dispose and TryDispose

|  |  |  |  |
| --- | --- | --- | --- |
| Type | Name | Params | Description |
| void | *Dispose* |  | Disposes the current vault. If the protected resource is *IDisposable*, it is disposed. Note that this method will throw an exception if it cannot obtain a lock in the time specified by the *DisposeTimeout* property. You should be sure that all threads that might be obtaining and releasing locks have stopped before calling. Alternatively, use *TryDispose*. |
| bool | *TryDispose* | [timeout : TimeSpan] | Try to dispose the vault for up to the time specified by timeout. Returns true for success, false for failure (indicating unable to obtain lock in time specified). |

Figure – Public Methods Common to All Vaults

#### Lock and SpinLock Methods

*Lock* and *SpinLock* perform the same function: they attempt to obtain the locked resource. *Lock* will put the thread to sleep for the *SleepInterval*[[27]](#footnote-27)in between failed attempts to obtain the resource. *SpinLock*, however, will actively loop until it obtains the resource or fails to do so. Both methods return whatever type of *LockedResourceObject[[28]](#footnote-28)* the Vault returns. The return value MUST be 1- declared inline[[29]](#footnote-29) and 2- be protected by a using statement or declaration *immediately* or the code will not compile.

The *Lock* and *SpinLock* methods are common to all vaults intended for direct consumption by users. These methods attempt to obtain a LockedResourceObject whereby access to the protected resource may be exclusively obtained. All these methods either succeed or throw an exception and the return value is always annotated with the *UsingMandatory* attribute which requires – as a condition of compilation – that the resource returned be protected by a *using* statement or declaration. The exceptions throwable by these methods are:

* + 1. *ObjectDisposedException*, indicating that the vault is being or has been disposed,
    2. *TimeoutException* indicating that the resource could not be obtained in the specified or fallback time window,
    3. *OperationCanceledException*, indicating that there was a request made via a CancellationToken to terminate the attempt prematurely, and
    4. *ArgumentOutOfRangeException* indicating that the timeout period specified was not positive.

Since one of the major difficulties in dealing with synchronization in large projects is the difficulty of reproducing and debugging deadlocks, the Vaults are designed to be difficult to deadlock.[[30]](#footnote-30) Thus, the parameterless overloads of *Lock* and *SpinLock* **do not attempt to lock forever, but instead imply specification of their** *DefaultTimeout* property as the timeout. If you desire an indefinite wait, use the overload that accepts a cancellation token and simply never propagate a cancellation request to it.

The overloads of these methods are summarized below:

|  |  |  |
| --- | --- | --- |
| ***Lock* and *SpinLock* Overloads** | | |
| **Overload Params** | **Description** | **Throws** |
| Parameterless | Attempt to obtain resource for the vault's default timeout period. | *ObjectDisposedException; TimeoutException;* |
| *TimeSpan* | Attempt to obtain resource for period specified by *TimeSpan* parameter | *ObjectDisposedException; ArgumentOutOfRangeException; TimeoutException;* |
| CancellationToken | Attempt to obtain resource until successful, or a cancelation request is propagated to the CancellationToken | *ObjectDisposedException; OperationCanceledException* |
| *TimeSpan; CancellationToken* | Attempt to obtain until the earliest of: a- successfully obtaining resource, b- period specified by *TimeSpan* parameter exceeded, c- cancellation request propagated to *CancellationToken* parameter | *ObjectDisposedException; ArgumentOutOfRangeException; TimeoutException; OperationCanceledException* |

Figure – Lock and Spinlock Overloads

### *BasicVault<T>*

The BasicVault is, by far, the easiest vault to use, but it comes with a significant limitation: it can only protect Vault-Safe resources. You will thus be limited to certain classes of value types and immutable reference types. If you acquaint yourself with the idioms of programming with immutable data structures, this is not as difficult: C#’s immutable collections[[31]](#footnote-31) have highly convenient methods that can create new collections with different values base on their current contents. The BasicVault can protect resources such as *DateTime*, *string, ulong*, as well as immutable collections of such values.

Because the BasicVault protects values that are intrinsically vault-safe, there are BasicVault supplies some convenience functions that other vaults cannot provide. These functions include:

* T CopyCurrentValue(TimeSpan timeout)[[32]](#footnote-32)
* (T value, bool success) TryCopyCurrentValue(TimeSpan timeout)[[33]](#footnote-33)
* void SetCurrentValue(TimeSpan timeout, T newValue)[[34]](#footnote-34)
* bool TrySetNewValue(TimeSpan timeout, T newValue)[[35]](#footnote-35)

### *MutableResourceVault<T>*

These vaults are designed to protect mutable resources and are far more restrictive than the *BasicVault*. The key to understanding **provably** *thread-safe interaction* with a mutable resource is to realize that all such interactions with the mutable resource must only read from or write to vault-safe objects. Consider updating a *StringBuilder* object, a paradigmatic example of a resource that is mutable and not vault-safe. **The data used to update the builder must itself be vault-safe.[[36]](#footnote-36)** Otherwise, mutable state accessible from outside the vault could cause a change in the state of the protected resource. This is the precise area where a reference-based, garbage collected languages like C# are at cross-purposes with thread-safety.[[37]](#footnote-37)

#### Factory Method

Special care must be taken when constructing a mutable resource vault. You must ensure that **only the mutable resource vault ever can see a reference to the value it protects**. To facilitate this, the factories for the mutable resource Vault require a delegate of type *Func<T>* as well as a positive *TimeSpan* defining the *DefaultTimeout* for the vault. You must ensure that the object constructed by the *Func<T>* -- as well as any non-vault-safe subparts of it are not accessible anywhere outside the delegate. The following example code shows correct and incorrect ways to accomplish this:



Figure – Correct and Incorrect Creation of MutableResourceVault

Care must also be taken to ensure that any non-vault-safe sub-objects of the constructed protected resource are not visible outside of the vault. The following example shows slightly more subtle correct and incorrect construction of a *MutableResourceVault*:



Figure – More Elaborate Correct and Incorrect Creation of MutableResourceVault

After correctly constructing a MutableResourceVault such that neither it nor any non-vault-safe sub-object thereof are accessible from outside the MutableResourceVault, the static analysis rules enable a high degree[[38]](#footnote-38) of confidence that no externally accessible mutable state will be mingled with the protected resource in the vault nor will any such mutable state be leaked out of the vault. The responsibility for *constructing* the object in a way that does not expose it or any non-vault-safe subpart thereof lies solely with the user.

#### Public Methods and Properties

The methods and properties exposed by the MutableResourceVault are like those exposed by BasicVault.[[39]](#footnote-39) The difference in usage lies not in the *MutableResourceVault* itself, but in the functionality and flexibility of its LockedResourceObject.

### *CustomizableMutableResourceVault<T>*

As shown below, the usage of the *MutableResourceVault*’s LockedResourceObjectlends itself to a manageable yet somewhat awkward and inconvenient syntax. For that reason, DotNetVault provides the *CustomizableMutableResourceVault* and examples showing how to create your own Vault Objects with their own LockedResourceObjects to which you can supply common operations in a more convenient syntax. The project itself provides a well-documented class called *StringBuilderVault* and its LockedResource object called *LockedStringBuilder* showing how to do this for your own classes and provide a very easy-to-use wrapper around the MutableResourceVault. It is hoped, as a future feature, to add code generation facilities to remove the boilerplate aspect of making these. That said, it is quite possible to use *MutableResourceVault* on its own and designing your own custom Vaults and LockedResource objects is a largely tedious mechanical task that is not difficult. If a vault protecting a non-vault-safe resource is to be used extensively, it is highly recommended that you take the time to customize the vault that protects it. Extension methods defined on the locked resource object also provide an option for improved readability.[[40]](#footnote-40)

## LockedResourceObjects In-Depth

### Common Functionality

All LockedResourceObjects are *ref structs[[41]](#footnote-41)* that can only be located on the stack. Such objects cannot be boxed, stored in static memory or made the non-static field of any type that is not itself a *ref struct*. Each has a Dispose method that returns the resource it guards to that vault from which it was obtained. When released for public consumption, the method returning these objects is always annotated with the *UsingMandatoryAttribute* which mandates that the immediate caller of the method immediately guard the return value with a *using* statement or declaration. This rule prevents failure to promptly Dispose the object, which is always a bug.

### *LockedVaultObject<TVault, [VaultSafeTypeParam] T>*

This type of LockedResourceObject is the type associated with the *BasicVault<T>*.[[42]](#footnote-42) *TVault* must be of type *BasicVault<T>*; *T* is the type of protected resource. The *VaultSafeTypeParamAttribute* informs the static analyzer that all type arguments matched with this parameter must be vault-safe. This object simply exposes the protected resource as a gettable and settable property. Since its type must be vault-safe, one need not worry about mingling it with unprotected resources or retaining a copy: if it is a reference type, the object to which it refers is immutable; if a value type, it is a deep copy or *effectively* a deep copy.[[43]](#footnote-43) Care should be taken with mutable structs: when accessing a struct through a property, you are accessing a copy of it. The following code sample demonstrates this care:



Figure – Care is Needed When Working With Mutable Structs

The output of executing such code should not be shocking to anyone familiar with Value Types and the .Net framework: accessing the Value through the getter *returns a copy, a* ***temporary*** *value*. The ubiquity of reference semantics, however, makes it easy to forget how this works.

OUTPUT:



Figure – Output Mutable Struct Demonstration

There are performance reasons that may make it desirable to be able to access the protected struct value by reference or constant reference. A future version of this software, with any analyzers needed to protect from abuse, may be released with this functionality.

### *LockedVaultMutableResource<TVault, TResource>*

This LockedResourceObject is associated with the *MutableResourceVault<T>*[[44]](#footnote-44)*.* Unlike the BasicVault’s LockedResourceObject, this LockedResourceObject comes with many restrictions, enforced by the static analyzer, to ensure that no mutable state from outside mingles with the protected resource and to ensure that no mutable state from the protected resource leaks outside of the vault. To facilitate this, all access to the resource is mediated through delegates annotated with attributes that are meaningful to the static analyzer. *TVault* must be of type *Vault<T>*.

#### LockedVaultMutableResource Delegates

The delegate declarations from the project are laid out:



Figure -- Special Delegates Used By LockedVaultMutableResource Objects to Prevent Leakage and Mingling of State

cont’d: 

Figure -- (cont’d) Special Delegates Used By LockedVaultMutableResource Objects to Prevent Leakage and Mingling of State

The static analyzer enforces the following rules on all these delegates:

* Their return value, if any, must be vault-safe.
* If an ancillary parameter is used, it must be vault-safe
* Anything captured in the closure assigned to the delegate must be vault-safe
* Other than the protected resource, anything that is read-from or written-to must be vault safe
* No reference to the protected resource or anything not vault-safe can be assigned to anything outside the delegate’s body.

*VaultQueries* are used to change no state in the mutable resource but to return information regarding the current state of the protected resource. The protected resource is passed to the delegate by constant reference. *VaultActions* return no values but are used to change the state of the protected resource. One ancillary vault-safe parameter is available. Only vault-safe captured values may be used. The protected resource is passed to the delegate by non-constant reference. *VaultMixedOperations* return values but also may change the state. The return value must be vault-safe. Any ancillary parameter used as well as anything captured in the closure must be vault-safe.

#### LockedVaultMutableResource Public Methods

The LockedVaultMutableResource exposes the following methods (with overloads allowing for an optional vault-safe ancillary parameter):

* *ExecuteQuery*– accepts a *VaultQuery* delegate, an optional vault-safe ancillary parameter and returns a vault-safe value based on the logic of the delegate. The protected resource is passed to the delegate by constant reference.
* *ExecuteAction*– accepts a *VaultAction* delegate, optional vault-safe ancillary parameter and changes the state of the protected resource as specified by the delegate. The protected resource is passed to the delegate by (non-constant) reference.
* *ExecuteMixedOperation*– accepts a *VaultMixedOperation* delegate, an optional vault-safe ancillary parameter, returns a vault-safe return value as specified by the logic of the delegate and may change the state of the protected resource. The protected resource is passed to the delegate by (non-constant) reference.

The best way to demonstrate the operation is to show examples of code with the output. For each example, assume that the protected resource is a *StringBuilder* that begins with the contents **“Hello, world!”**.



Figure -- VaultQuery Demonstration

Figure -- VaultQuery Demo Output

Performing DemonstrateQueries...

Contents: Hello, world!.

Length of contents (content length: [13]) plus [ancillaryValue] of [7]: 20

Char at idx [1] (current val: [e]) made upper case: [E].



Figure ­-- VaultAction Demonstration

Figure ­-- VaultAction Demo Output

Performing DemonstrateActions...

Reversed Upper/Lower res: hELLO, WORLD!

Made chars at idx divisble by 3 q: [qELqO,qWOqLDq]



Figure – VaultMixedOperation Demonstration

Figure -- VaultMixedOperation Demo Output

Performing DemonstrateMixedOperations...

New value: [Hello it's magic oooh oooh! , world!], Index: [4].

### 

As can be seen, interacting with protected mutable resources requires more care and requires understanding of expression and statement lambdas. As mentioned above, facilities exist to create customized, more convenient LockedResourceObjects and code generation utilities are planned for future released.[[45]](#footnote-45) A shorter path to a more convenient syntax is also available: passing the locked resource to extension methods (by reference). The following example shows how extension methods can be used to simplify frequently used syntax:



Figure – Demonstration of Extension Methods to Simplify Usage

Figure -- Output of Extension Method Demo

Performing DemonstrateUseOfExtensionMethodsToSimplify...

Contents: Hello, world!

First char: H

Changed to uppercase 'E': HEllo, world!

# Static Analyzer Rules

The static analyzer facilitates isolation of protected resources and prevents inadvertent failure to Dispose a LockedResourceObject.

## DotNetVault\_UsingMandatory

### Summary

This rule requires that the caller to any method that annotates its return type with the *UsingMandatoryAttribute* protect that return value with a using statement or declaration immediately. Failure to do so will cause a compilation error. LockedResourceObjects represent a lock that has been obtained on a protected resource that more than one thread wishes to access. Allowing it to be held open for longer than its lexical scope can starve out other threads and potentially result in the inability of any other thread to ever be able to access the locked resource again. If no timeout is used, it would result in deadlock. Use of the using statement is simple syntactically and guarantees that the lock will be disposed (and the resource returned to the vault) no later than the end of the scope in which it is received from the method returning the value.

### Known Bug (#50)

There exists a known bug whereby a user can predeclare a *LockedResourceObject*, assign the value later in a using statement, then access the disposed value after the end of the scope of the using statement. This behavior is not supported, and the bug will be fixed by making it an error to fail to declare and assign the locked resource value inline.



Figure -- Bug #50 Demonstration

## DotNetVault\_VaultSafe

This rule enforces that any type annotated with the vault-safe attribute[[46]](#footnote-46), without its *OnFaith* property set to true, must strictly comply with all requirements imposed on vault-safe types. If you use the *OnFaith* property on a type it will be deemed vault-safe without analysis. This allows types that partially employ their own thread-safety mechanisms to be considered vault-safe by the analyzer. You must be careful because any changes to the type will not cause any compiler errors or warnings despite their affecting the vault-safety of the type.

## DotNetVault\_VsDelegateCapture

This rule ensures that none of the LockedVaultMutableResource delegates[[47]](#footnote-47) read or write to any captured or static variable that is not vault-safe except for the protected resource itself. This prevents any externally accessible mutable state from leaking outside or creeping into it. Within these delegates, it is an error to read from or write to anything save vault-safe data and the protected resource itself.

1. DotNetVault\_VsTypeParams,

DotNetVault\_VsTypeParams\_MethodInvoke, DotNetVault\_VsTypeParams\_ObjectCreate, and

DotNetVault\_VsTypeParams\_DelegateCreate

These rules enforce that a whenever a type argument is substituted for a type parameter annotated with the *VsTypeParamsAttribute* that type is vault-safe. Essentially, it introduces a new type-constraint: that of vault-safety. The first rule enforces this in type declarations. The second whenever type arguments are provided to a generic method, the third at the instantiation of objects and the final when generic delegates are created. These rules are primarily applicable in the contexed of LockedVaultMutableResource delegates.[[48]](#footnote-48)

1. DotNetVault\_NotVsProtectable

Frequently a type’s API will expect to receive collections of objects rather than individual objects, the most general of which is *IEnumerable<T>*. Neither *IEnumerable<T>* nor the most used collections in .NET’s base class library are vault-safe. While, ideally, Immutable Collections (with vault-safe type arguments) would be used for updating protected resources, this could produce inconvenience and perhaps unacceptable performance on occasion.

To mitigate this problem, DotNetVault provides convenience types designed to wrap generic collections whose types are vault-safe with minimal performance overhead:

|  |  |  |
| --- | --- | --- |
| Vault-Safe Convenience Wrappers | | |
| Declared Type | Use Case | Description |
| public struct ArrayEnumeratorWrapper<[VaultSafeTypeParam] T> : IEnumerator<T> | Takes an array of a Vault-Safe Type and provides a struct Enumerator to iterate it. Access to the underlying mutable collection is prevented. | Functions like any enumerator. As a struct enumerator, no unnecessary indirection or GC Pressure is introduced. |
| public struct KvpEnumerator<[VaultSafeTypeParam] TKey, [VaultSafeTypeParam] TValue> : IEnumerator<KeyValuePair<TKey, TValue>> | Takes a collection of KeyValuePair (where both TKey and TValue are vault-safe) and provides a struct enumerator to iterate it. Access to underlying mutable collection is prevented. | Functions like any enumerator. As a struct enumerator, no unnecessary indirection or GC Pressure is introduced. |
| public struct StandardEnumerator<[VaultSafeTypeParam] T> : IEnumerator<T> | Take any IEnumerator<T> of a vault-safe type T and wrap it in a struct enumerator that prevents access to any underlying mutable collection | Functions like any enumerator. As a struct enumerator, no *additional* indirection or GC Pressure is introduced. |
| public struct StructEnumeratorWrapper<TWrappedEnumerator, [VaultSafeTypeParam] TItem> : IEnumerator<TItem> where TWrappedEnumerator : struct, IEnumerator<TItem> | Take any existing struct enumerator of a vault-safe type T and return an enumerator that will prevent access to the underling mutable type. | Functions like any enumerator. As a struct enumerator, no *additional* indirection or GC Pressure is introduced. |
| public sealed class VsArrayWrapper<[VaultSafeTypeParam] T> : IVsArrayWrapper<T> | Take any existing array of vault-safe type T, and wrap it in a class that allows readonly access to the array but prevents changes | Functions like an *IReadOnlyList*<T>. Some overhead and GC pressure introduced, but minimal compared to overhead introduced by copying a large collection. |
| public sealed class VsEnumerableWrapper<[VaultSafeTypeParam] T> : IVsEnumerableWrapper<T> | Take any existing collection of vault-safe values (IEnumerable<T>) and wrap it in a collection that prevents access to any underlying mutable collection | Functions like any *IEnumerable*<T>. Some overhead and GC pressure introduced, but minimal compared to overhead introduced by copying a large collection. |
| public sealed class VsListWrapper<[VaultSafeTypeParam] T> : IVsListWrapper<T> | Take any existing collection of vault safe values of type List<T> and wrap it in a read-only collection that prevents access to any underlying mutable collection | Functions like an IReadOnlyList<T>. Some overhead and GC pressure introduced, but minimal compared to overhead introduced by copying a large collection. |

Figure -- Vault-Safe Convenience Wrappers

These types are all annotated with the *VaultSafeAttribute* with its *OnFaith* property set to true but are also annotated with the *NotVsProtectableAttribute[[49]](#footnote-49).* The purpose of the wrapper is to permit vault-safe access to values of these types in *LockedVaultMutableResource* object’s delegates without fear of permitting access to the underlying mutable collection therein. *The are* ***not****, however,* thread-safe in any way. If you use them, you must be sure that no other threads have write access to these objects while they are being processed by the delegate. Because these objects do not provide real thread-safety or vault-safety (but instead a temporary façade of such safety), the *NotVsProtectableAttribute* forbids them from being a part of any protected resource in any vault. Their sole purpose is conveniently to prevent mingling or leaking mutable state with the protected resource.

This rule detects and forbids use of these objects as a protected resource in a vault or being a part, in any way, of such a protected resource. If you wish to protect a collection of vault-safe types, either use the BasicVault to protect an immutable collection thereof or use the MutableResourceVault and write extension methods or a customized vault permitting updates via immutable collections or these convenience wrappers.

The following shows the intended use-case for such wrappers:



Figure – Usage of Vs Convenience Wrappers

Begin showing wrapper usage.

Printing final contents:

{1, 2, 3, 4, 5, 6, 7, 8}

Done showing wrapper usage.

Figure ­­-- Usage Wrapper Demo Output

# Attributes

## VaultSafeAttribute

The vault-safe attribute when placed on a class or struct indicates to the static analyzer that the type so-annotated is vault-safe. A default-constructed instance of this attribute, or one with the first constructor parameter evaluating to *false*, will have the static analyzer verify that the type is provably vault-safe.[[50]](#footnote-50) If constructed with its first constructor parameter evaluating to *true*, the static analyzer shall assume, without further analysis, it to be vault-safe. Note that types that comply with the *unmanaged* type constraint are deemed vault-safe regardless of the presence or absence of this attribute.

There is also included herewith a whitelist file called “VaultSafeWhiteList.txt” with the names of types should be deemed vault-safe by the analyzer without analysis. The following shows the by-default contents of this whitelist:

String; System. Uri;

Figure -- Contents of Whitelist.txt

There is a second whitelist file for conditionally vault-safe generic types called “condit\_generic\_whitelist.txt”. These types will be deemed vault-safe without analysis if and only if all their generic type arguments are vault-safe. This is included to allow the use of immutable collections (whose generic type arguments are all vault-safe) as vault-safe types. If you design any generic immutable collections, they will be deemed vault-safe whenever all their type arguments are vault-safe. The starting contents of the file are shown below:

## UsingMandatoryAttribute

Figure -- Contents of condit\_generic\_whitelist.txt

System.Collections.Immutable.ImmutableHashSet`1;

System.Collections.Immutable.ImmutableList`1+Enumerator;

System.Collections.Immutable.ImmutableQueue`1+Enumerator;

System.Collections.Immutable.ImmutableArray`1+Enumerator;

System.Collections.Immutable.ImmutableHashSet`1+Enumerator;

System.Collections.Immutable.ImmutableSortedSet`1;

System.Collections.Immutable.ImmutableDictionary`2+Enumerator;

System.Collections.Immutable.ImmutableSortedDictionary`2+Enumerator;

System.Collections.Immutable.ImmutableList`1;

System.Collections.Immutable.ImmutableStack`1+Enumerator;

System.Collections.Generic.KeyValuePair`2;

System.Collections.Immutable.ImmutableQueue`1;

System.Collections.Immutable.ImmutableArray`1;

System.Collections.Immutable.ImmutableSortedSet`1+Enumerator;

System.Collections.Immutable.ImmutableStack`1;

System.Collections.Immutable.ImmutableDictionary`2;

System.Collections.Immutable.ImmutableSortedDictionary`2;

This attribute may be used to annotate the return value of a method. When so annotated, it becomes a compiler error for the receiving attribute to fail to declare and assign the variable in the context of a *using* statement or declaration.[[51]](#footnote-51)

## VaultSafeTypeParamAttribute

This attribute can be applied to type parameters. It effectively serves as a type constraint, enforced by the static analyzer, that any type argument substituted for the parameter must be vault-safe. It is considered an error for such a type not to be vault-safe. This may attribute may annotate type parameters wherever they are declared: on generic types, generic methods and generic delegates.

## NoNonVsCaptureAttribute

This attribute annotates delegates causing the static analyzer to ensure that all variables, not explicitly passed to it as method arguments, accessed by the delegate are vault-safe. It is an error to capture or otherwise read to or write from any variable from within the delegate unless the variable is:

1. vault-safe or
2. passed as a type argument

The LockedVaultMutableResource delegates[[52]](#footnote-52) are all annotated with this attribute. All type parameters passed, except the first, which is the protected resource, are annotated with the *VaultSafeTypeParamAttribute*. This combination of attributes prevents any non-vault-safe variables, other than the protected resource, from being read from or written to in the delegate.

## NotVsProtectableAttribute

This attribute may be applied to a type. Any type – even a vault-safe type – so annotated cannot be used as a protected resource or part of a protected resource in any vault. If the type is otherwise vault-safe, it will be considered vault-safe only for purposes of the rules at § [5.b-d](#_DotNetVault_VaultSafe), *supra*. This attribute is used primarily to allow for convenience wrappers around mutable collections of vault-safe type to be used in LockedVaultMutableResource delegates.

# Licensing

## Software License

The DotNetVault software library and source-code are the intellectual property of CJM Screws, LLC, a Maryland Limited Liability Company. CJM Screws, LLC licenses the software and source code to you under the MIT License as follows:

MIT License

Copyright (c) 2019 CJM Screws, LLC

Permission is hereby granted, free of charge, to any person obtaining a copy

of this software and associated documentation files (the "Software"), to deal

in the Software without restriction, including without limitation the rights

to use, copy, modify, merge, publish, distribute, sublicense, and/or sell

copies of the Software, and to permit persons to whom the Software is

furnished to do so, subject to the following conditions:

The above copyright notice and this permission notice shall be included in all

copies or substantial portions of the Software.

THE SOFTWARE IS PROVIDED "AS IS", WITHOUT WARRANTY OF ANY KIND, EXPRESS OR

IMPLIED, INCLUDING BUT NOT LIMITED TO THE WARRANTIES OF MERCHANTABILITY,

FITNESS FOR A PARTICULAR PURPOSE AND NONINFRINGEMENT. IN NO EVENT SHALL THE

AUTHORS OR COPYRIGHT HOLDERS BE LIABLE FOR ANY CLAIM, DAMAGES OR OTHER

LIABILITY, WHETHER IN AN ACTION OF CONTRACT, TORT OR OTHERWISE, ARISING FROM,

OUT OF OR IN CONNECTION WITH THE SOFTWARE OR THE USE OR OTHER DEALINGS IN THE

SOFTWARE.

## Documentation License

This document is the intellectual property of CJM Screws, LLC, a Maryland Limited Liability Company. CJM Screws, LLC hereby grants you the right to use, copy and distribute this document in unaltered form without limitation provided any persons to whom the documentation is distributed agree to be bound by the terms of this license including the following:

1. The documentation may not be altered in any material way and
2. THE DOCUMENTATION IS PROVIDED "AS IS", WITHOUT WARRANTY OF ANY KIND, EXPRESS OR IMPLIED, INCLUDING BUT NOT LIMITED TO THE WARRANTIES OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE AND NONINFRINGEMENT. IN NO EVENT SHALL THE AUTHORS OR COPYRIGHT HOLDERS BE LIABLE FOR ANY CLAIM, DAMAGES OR OTHER LIABILITY, WHETHER IN AN ACTION OF CONTRACT, TORT OR OTHERWISE, ARISING FROM, OUT OF OR IN CONNECTION WITH THE DOCUMENTATION OR THE USE OR OTHER DEALINGS IN THE DOCUMENTATION.
3. *Author Contact Information*

The author of the documentation and software is Christopher P. Susie, a member of CJM Screws, LLC. He can be reached by email at [cpsusie@hotmail.com](mailto:cpsusie@hotmail.com) or by mail at

CJM Screws, LLC

ATTN: Christopher P. Susie

3705 Ridgecroft Road

Baltimore, Maryland 21206

U.S.A.

1. <https://docs.microsoft.com/en-us/dotnet/csharp/language-reference/keywords/lock-statement> [↑](#footnote-ref-1)
2. RAII stands for Resource Acquisition is Initiation. It is the primary method used in modern C++ to manage resource lifetime and is also the method employed by the Rust programming language. It is a rather poorly thought-out name. Its principle is simple: when an object is initialized it acquires all its resources and establishes all invariants – making it impossible for an uninitialized or invalid object to exist. Equally, and perhaps more importantly, when an object’s lifetime ends (typically bound to scope), at that moment it releases any resources it acquired. <https://en.cppreference.com/w/cpp/language/raii>; <http://www.stroustrup.com/bs_faq2.html#finally> [↑](#footnote-ref-2)
3. v. *supra,* #1. [↑](#footnote-ref-3)
4. <https://docs.microsoft.com/en-us/dotnet/api/system.threading.mutex?view=netframework-4.8> [↑](#footnote-ref-4)
5. <https://docs.microsoft.com/en-us/dotnet/api/system.threading.spinwait?view=netframework-4.8> [↑](#footnote-ref-5)
6. <https://docs.microsoft.com/en-us/dotnet/api/system.threading.spinlock?view=netframework-4.8> [↑](#footnote-ref-6)
7. <https://docs.microsoft.com/en-us/dotnet/api/system.threading.interlocked?view=netframework-4.8> [↑](#footnote-ref-7)
8. v. #1, *supra.* [↑](#footnote-ref-8)
9. v. Figure 1, *supra.* [↑](#footnote-ref-9)
10. v. #1, *supra.* [↑](#footnote-ref-10)
11. <https://en.wikipedia.org/wiki/Reentrant_mutex#cite_note-2> [↑](#footnote-ref-11)
12. *Id.*  [↑](#footnote-ref-12)
13. Would all developers on your team think of this? Of those that would, will they realize they need to do it even when under other pressures? Is this type of thing really a good idea anyway? [↑](#footnote-ref-13)
14. The Liskov substitution principle requires a derived class to be a drop-in substitute for its base classes. Here, however, extra code is needed to use the thread-safe list correctly meaning that it is not a drop-in substitute for the list it replaces. *See,* *e.g.*, <https://en.wikipedia.org/wiki/Liskov_substitution_principle> [↑](#footnote-ref-14)
15. *See e.g.* <https://doc.rust-lang.org/book/ch16-03-shared-state.html#using-mutexes-to-allow-access-to-data-from-one-thread-at-a-time> (explaining that Rust’s approach to shared mutable state attacks the “shared” problematic aspect by preventing access to the resource without first obtaining the lock); <https://github.com/facebook/folly/blob/master/folly/docs/Synchronized.md> (reiterating that conventional synchronization is error prone because reliant solely on conventions and explaining that Folly’s Synchronized library solves this problem by inextricably linking access to the data with obtaining the lock). [↑](#footnote-ref-15)
16. In pre-2011 C++ idioms, it was common to allocate objects on the heap with new and delete them with delete. These objects were often accessed by pointer. The difficulty with this was that you had to be sure the underlying object was deleted *exactly once* and, *once deleted*, never accessed again (accounting for the possibility of early return or exceptions). Sane code thus required careful circumspection. Modern C++ eschews direct heap allocation in most instances and prefers allocating concrete objects on the stack where possible. Concrete objects on the stack are automatically cleaned up when they go out of scope. When “owning” pointers are used, the use of smart pointers such as std::unique\_ptr or std::shared\_ptr are preferred. These smart pointers manage deletion of the object automatically. [↑](#footnote-ref-16)
17. Instead, it is better to realize that the value of a value type is stored directly in the variable (or array) that represents it. Regular structs can reside in three places: for local variables: on the stack, for struct members of reference types, on the heap, and for static variables, in static memory. For structs that are elements of an array of that struct type, the *actual value* of the struct is stored on the heap, contiguously with other elements of the array. For structs that have been boxed into a reference type, the boxed struct resides on the heap. If such a boxed struct is an element in an array of objects or interfaces, the array contains references to the elements (some or all of which may be structs or classes), not the value of the objects themselves. In other words, in *TimeSpan[],* the values of the *TimeSpans* are stored **on the heap,** **contiguously, directly in the array**. In *IEquatable<TimeSpan>[]*, each element in the array is a reference to a boxed *TimeSpan* or some other object that implements this interface; only references to the elements are contiguous – the values themselves are scattered. [↑](#footnote-ref-17)
18. Boxing occurs to allow structs to implement interfaces and to expose the methods common to System.Object even if it does not override them itself. Simply being stored as a struct field in a class or as an item in array of structs is not boxing. Also, storing structs in a generic collection does *not* cause boxing. [↑](#footnote-ref-18)
19. <https://kalapos.net/Blog/ShowPost/DotNetConceptOfTheWeek16-RefStruct>; <https://blogs.msdn.microsoft.com/mazhou/2018/03/02/c-7-series-part-9-ref-structs/> . [↑](#footnote-ref-19)
20. <https://stu.dev/csharp8-doing-unsupported-things/> [↑](#footnote-ref-20)
21. “*unmanaged*” is a type constraint introduced in C# 7.3. According to the standard, a type is unmanaged if “is not a reference type and doesn't contain reference type fields at any level of nesting.” <https://docs.microsoft.com/en-us/dotnet/csharp/language-reference/proposals/csharp-7.3/blittable> . Types parameters that meet the requirements of this constraint can have their address assigned to a pointer inside *unsafe* genericmethods. [↑](#footnote-ref-21)
22. v. §§ [6.a](#_VaultSafeAttribute), [5.b](#_DotNetVault_VaultSafe), *infra*. [↑](#footnote-ref-22)
23. [*Id*](#_VaultSafeAttribute)*.* [↑](#footnote-ref-23)
24. *See* Concept of Vault-Safety*, supra* at a.  [↑](#footnote-ref-24)
25. You *may* use a *MutableResourceVault* to protect a vault-safe object as well. It may be advisable to use it if the resource you are protecting is under active development and might lose its Vault-Safe status later. If the resource is unlikely to need to become non-Vault-Safe, the BasicVault should be used because its locked resource is less restrictive and easier to work with. [↑](#footnote-ref-25)
26. See [§ 1.c.i](#_C#_8’s_Disposable) , *supra.* [↑](#footnote-ref-26)
27. v. [§ 4.c.3.1](#_Public_Read_Only), *supra*. This period should be *very* short. It currently defaults to 10 milliseconds. [↑](#footnote-ref-27)
28. v. [§ 4.b.2](#_LockedResources), *supra*. [↑](#footnote-ref-28)
29. There exists a bug in the analyzer (v. [§ 5.a.2](#_Known_Bug_(#50)), *infra*) that permits the locked resource to be pre-declared. This is not supported and the fix for Bug #50 will make out-of-line declarations a compiler error. Always declare the locked resource object inline. [↑](#footnote-ref-29)
30. The desire to make deadlock as difficult as possible is not *always* desirable as it can waste CPU time, which can be a problem when deployed at scale on a server under heavy load. A future version may provide a mutex and condition variable based version of the vaults which will not suffer from this drawback (but, of necessity, will be easier to deadlock). [↑](#footnote-ref-30)
31. These are contained in the *System.Collections.Immutable* namespace. They are included as part of .NET Core and are available as a NuGet package for the .NET framework. DotNetVault’s static analyzer (and Roslyn analyzers generally) make extensive use of them to allow concurrent execution of analysis without fear of corrupting mutable state. [↑](#footnote-ref-31)
32. Copies the current value protected by the vault. Will throw *ArgumentOutOfRangeException, TimeoutException and ObjectDisposedException*. Returns a copy of the protected resource. [↑](#footnote-ref-32)
33. Tries to copy the value currently protected by the vault. Will return (true, value) on success or (false, undefined) on timeout. Will throw *ObjectDisposedException* and *ArgumentOutOfRangeException*. [↑](#footnote-ref-33)
34. Replaces the value currently protected by the vault with a specified new value. Will throw *ArgumentOutOfRangeException*, *TimeoutException* and *ObjectDisposedException*. [↑](#footnote-ref-34)
35. Tries to replace the value currently protected by the vault with a specified new value. Will return true for success and false for a timeout failure. Will throw *ArgumentOutOfRangeException* and *ObjectDisposedException*. [↑](#footnote-ref-35)
36. One would not pass a *StringBuilder* or a *List<char>* to append to the end of the protected string builder. Instead one would pass a *string* (Vault-Safe) or an *ImmutableArray<char>* or *ImmutableList<char>* to the *StringBuilder*. In this way we can be sure that no subsequent external change to the objects referred to by these parameters will not affect the protected resource. [↑](#footnote-ref-36)
37. In a language like C++ or, especially, Rust where value semantics are used ubiquitously and objects can be, frequently are, and should be – wherever possible – on the stack, there is little danger of concurrent access. When passed by value, they are deep copies; when moved, the resource is no longer accessible from the original variable. In C++ and *unsafe* blocks in Rust this can be defeated by pointers, references and reference-like view objects (e.g. [std::string\_view](https://en.cppreference.com/w/cpp/string/basic_string_view) and [std::span](https://en.cppreference.com/w/cpp/container/span)). In C#, where most objects, by design, are on the heap and accessed indirectly, mutable state almost begs to be shared. [↑](#footnote-ref-37)
38. Although I have performed testing and believe that it will be difficult to mingle mutable state with the protected object or to allow mutable state to leak from the vault, I cannot guarantee it. This project is released, under the MIT license, with no warranties, not even the warranty of merchantability or fitness for any particular purpose. See the license agreement for more legal details. I welcome any reports of ways found that circumvent the protection afforded by this project’s static analyzer and welcome proposed fixes and new static analysis rules to make the protections provided by this project stronger. [↑](#footnote-ref-38)
39. v. [§ 4.c.3](#_Functionality_Common_to), *supra*. [↑](#footnote-ref-39)
40. v. Figure 21, *infra*. [↑](#footnote-ref-40)
41. v. #19, *supra*. [↑](#footnote-ref-41)
42. v. § [4.c.ii](#_BasicVault<T>), *supra*. [↑](#footnote-ref-42)
43. v. § [4.a](#_Concept_of_Vault-Safety), *supra*. [↑](#footnote-ref-43)
44. v. § [4.c.iii](#_MutableResourceVault<T>), *supra*. [↑](#footnote-ref-44)
45. v. § [4.c.iv](#_CustomizableMutableResourceVault<T>), *supra*. [↑](#footnote-ref-45)
46. v. § [4.6.a](#_VaultSafeAttribute), *infra*. [↑](#footnote-ref-46)
47. v. § [4.d.iii.1](#_LockedVaultMutableResource_Delegate), *supra*. [↑](#footnote-ref-47)
48. *Id*. [↑](#footnote-ref-48)
49. v. § [6.e](#_NotVsProtectableAttribute), *infra*. [↑](#footnote-ref-49)
50. For the characteristics of vault-safe types see § [4.a](#_Concept_of_Vault-Safety), *supra*. [↑](#footnote-ref-50)
51. v. § [5.a](#_DotNetVault_UsingMandatory), *supra*. [↑](#footnote-ref-51)
52. v. § [4.d.iii.1](#_LockedVaultMutableResource_Delegate), *supra*. [↑](#footnote-ref-52)